EECS2011 Fundamentals of Data Structures
(Winter 2022)

Q\&A - Week 4 Lecture

Wednesday, February 9

Announcements

- Written Test 1 due next Monday or Tuesday
- Revised start time of Written Test 1
- Example questions for Written Test 1 released
- Assignment 1 (on SLLs) due next Tuesday
- Lecture W5 postponed until next Wednesday


For some of the methods we do 1 block of error checking like this: methodl throws Exception

```
if () {
    throw exception
    }
else{
    //rest of the code
        //
    }
```

Sometimes I exclude the "else\{\}" o save 1 indentation level and
to minimize confusion when having multiple loose curly brackets at the end.
Since the code below the 'if block' can only be reached if there is no error anyways.
if () \{
throw exception
\}
// rest of the code

Question: Is this way of doing it bad programming practice?

Hi professor,
In the "addAt" method, if we want to keep track of the tail, should we add a case in the algorithm for that?
My thinking is to have a special case if the size of the list equals the index to be added.
In that case we simply use addLast. ie.:
if(this.size $==$ $\square$ ( i $^{2}$ ) \{ this.addLast(new Node(e,null); \} Else the algorithm runs as normal.
This allows us to keep track of the tail
and makes one special case faster. Else the algorithm runs as normal.
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this.addLast(new Node(enul);

size

$\left.\operatorname{addAt}(3), ~ " a \ln ^{\prime \prime}\right)$;


At 17:55 - in removeFirst(), we do
$\checkmark$ daltead
clanlode = head
class SLL \{ $\qquad$ Java
head $=$ head.getNext() $\cup$ Node dulled $\overline{\text { manalloc. }}$


Why do we have to do line + line 3 for the previous node O to be Garbage Collected ?
Wouldn't it be garbage collected if we only wrote line 2?

- since the variable head is no longer pointing to the previous node, and no other node or variable has a reference to the previous node 0 either?


Node dall-lead= head; head $=\frac{\text { head. get Next () ; }}{}$
dollar

$$
\begin{aligned}
& \text { int } \bar{i}=23 \overline{3} \\
& \text { int } \bar{J}=46 ; \\
& \tau=\mathrm{J} ;
\end{aligned}
$$

## Hi Professor, <br> VemoveAt( $\overline{\text { ant }} \overline{\text { I }}$ )

When we remove a node from a singly linked list,
despite disconnecting it from the previous or the next nodes whichever is applicable, I believe it will still occupy space in the memory.

How can we completely remove it so it does not occupy any space in the memory? I am thinking about a very large linked list and if we remove multiple nodes and they still occupy memory spaces, I think that would be a big problem. Thank you.


As a developer, when writing a method we might
use the parameter to call some methods defined in that parameter's static type.

For example, if our method is $m$ (Person $p$ ), when developing we might need to call the method p.get BMI() because we know that parameter $p$ has this method defined.

But when we set the method's parameter type generic $m(E p)$, how can we call a method (e.g., p.getBMI()) given that we know that we have to call that method to be able to write the method that we are currently writing?

What about a case where an argument with static type Account is passed, how can we avoid this?

Thanks

Constrained Genericity

private Element; private Node<E> next;
public double getBMIO \{ return this.element getBMI();
\}

public double getTuition() \{ double tuition; if(element instanceof Student) \{ tuition = ((Student) element).getTuition(); \} else \{ tuition $=-1$; \}
\}
public class Person \{ public double getBMI() \{ return 0;
public class Student extends Person \{ public double getTuition() \{ return 0; \}
\}


